



TraceMesh: Scalable and Streaming Sampling for Distributed Traces

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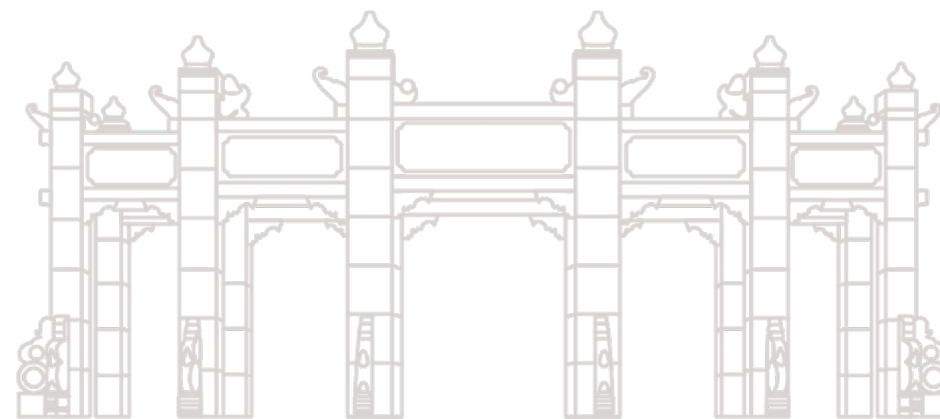
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Outline

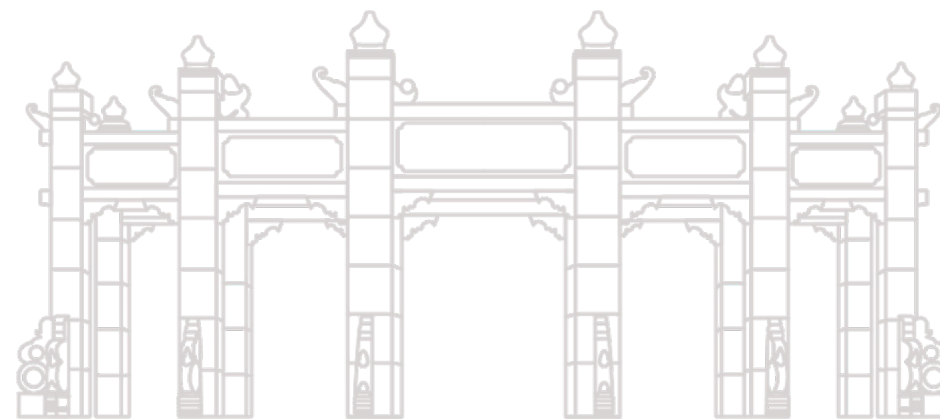


- **Introduction & Motivation**
- **TraceMesh Overview & Design**
- **Evaluation**
- **Conclusion**

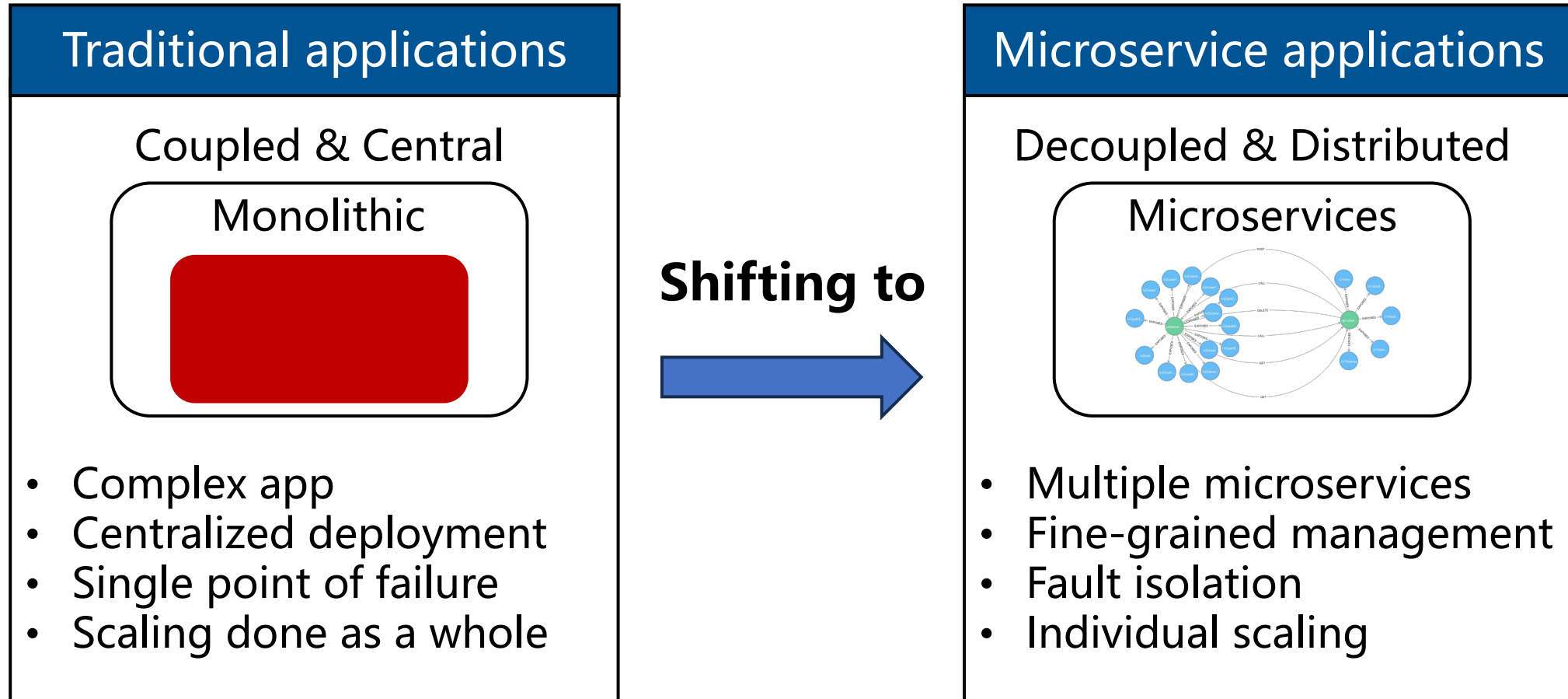




Introduction & Motivation

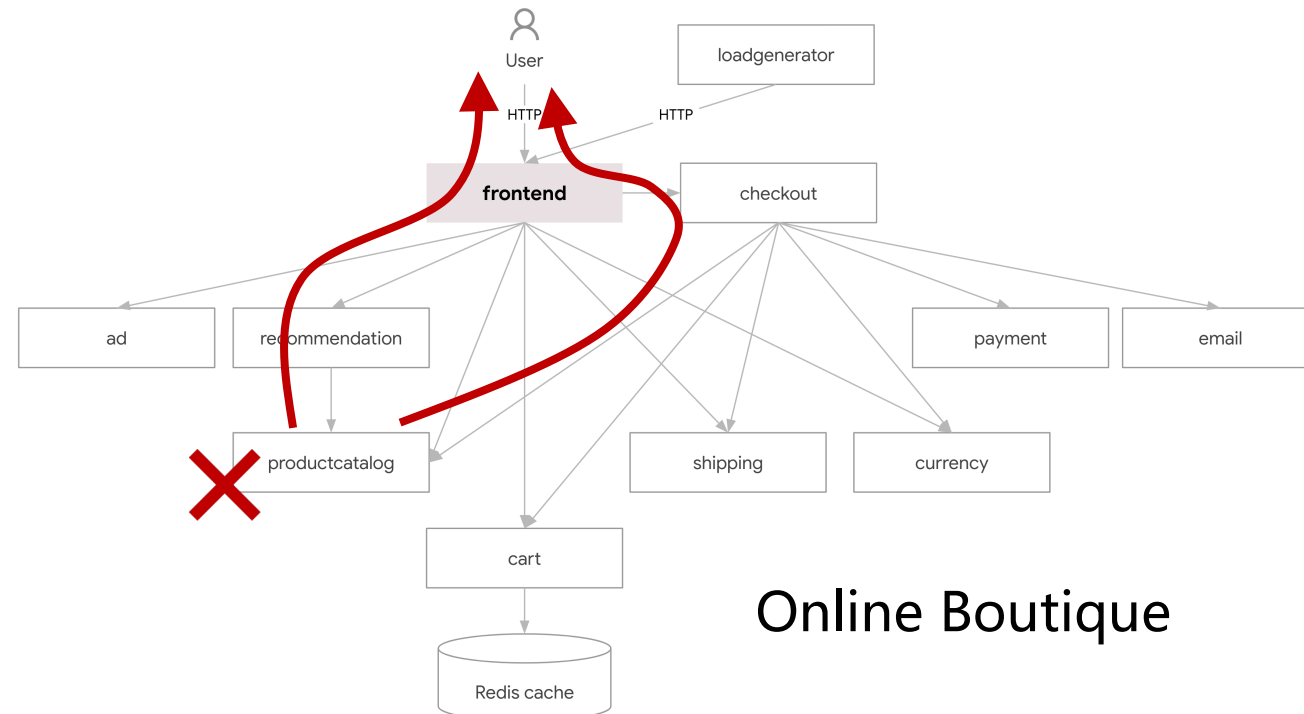


Shifting to Microservices



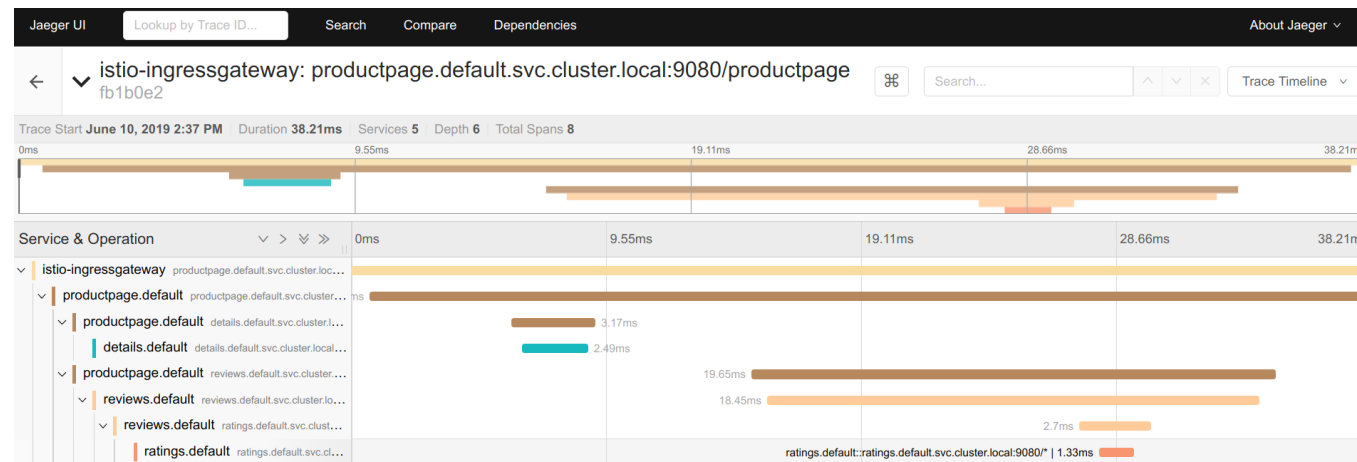
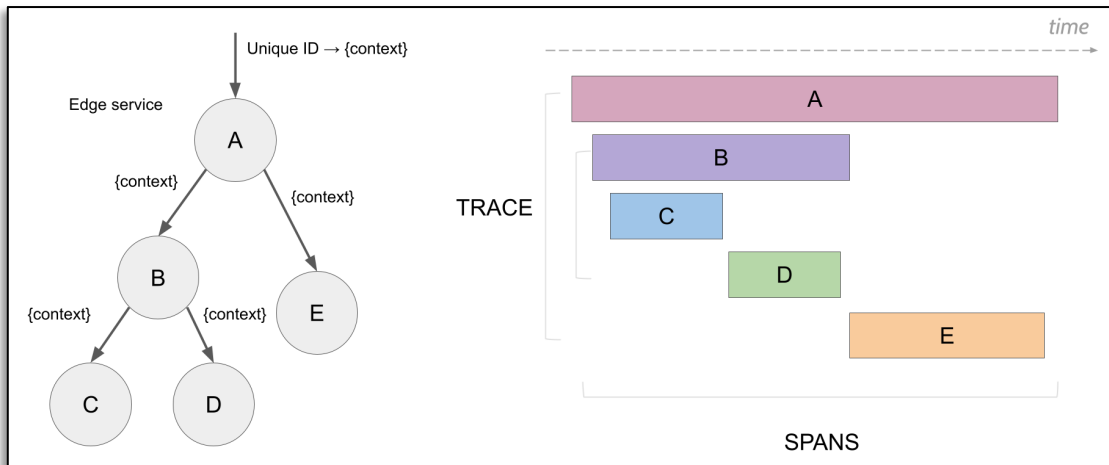
Complex Reliability Management

- E2E behavior can be affected by any microservice
- Failures cascading effect
- Symptoms and root causes can be far apart



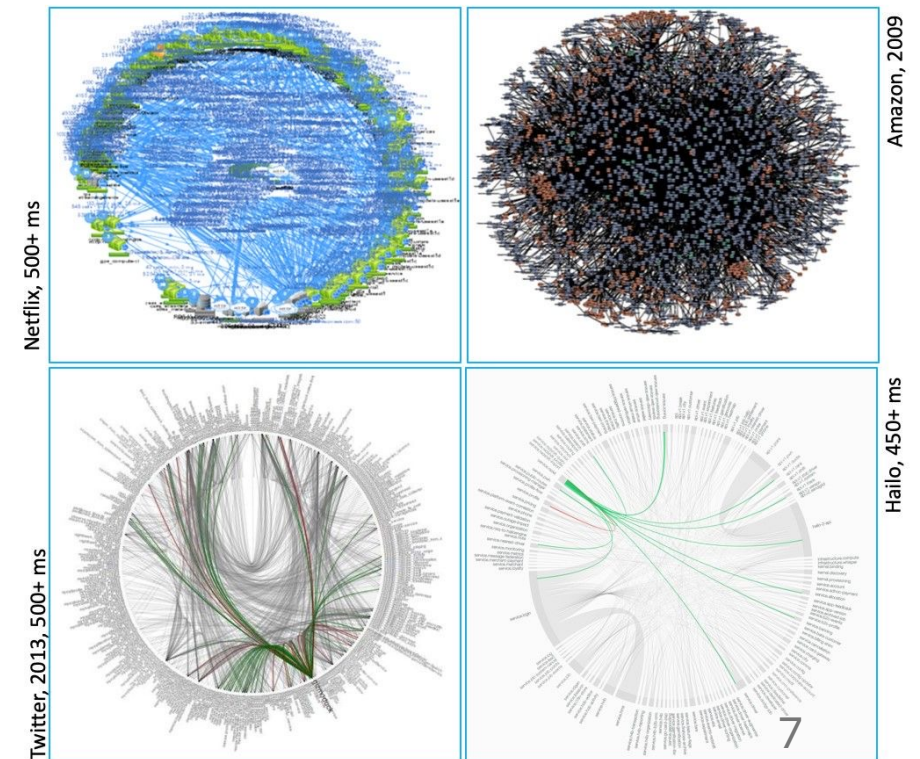
Distributed Tracing

- A directed tree recording executions across microservices
- Trace events: time, operations, messages, attributes
- Critical for performance monitoring, root cause analysis, etc.



Trace Data are Large in Size

- Complex microservices systems
 - Uber is composed of thousands of microservices
 - WeChat system hosts more than 3,000 services
- Huge volume of trace data
 - E.g., Google generates about 1,000 TB of raw traces on a daily basis
 - **Significant overhead to trace generation, collection, and ingestion**



Most Trace Data are Useless

- Most traces are recurring and useless
 - Cloud services often guarantee more than 99% uptime in their SLAs
 - Only edge-case traces are interesting and **should be sampled**
 - **By definition, these traces are rare**

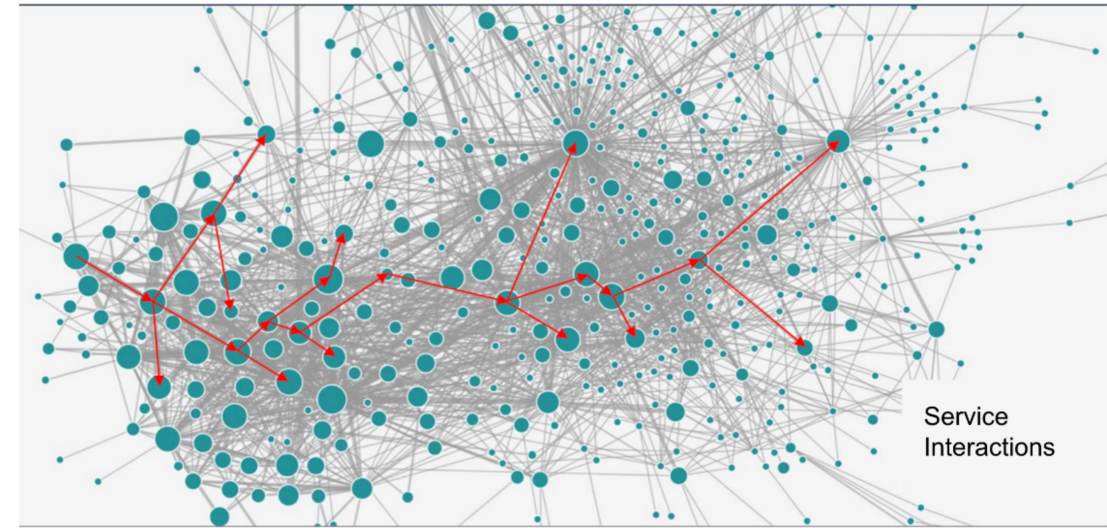
AMAZON ELASTIC COMPUTE CLOUD (EC2) (INSTANCE LEVEL)	
Monthly Uptime Percentage	Service Credit Percentage
Less than 99.5% but equal to or greater than 99.0%	10%
Less than 99.0% but equal to or greater than 95.0%	30%
Less than 95.0%	100%
Amazon Compute Full SLA	
Compute Containers	

Trace Sampling

- Head-based sampling
 - Decision made **prior** to request execution
 - Uniformly sample traces at random with a sampling rate, e.g., 1%
 - Could result in common-case traces
- Tail-based sampling
 - Decision made **after** request execution
 - Biased sampling based on latency, HTTP status code, etc.
 - Learning-based approaches to identify uncommon traces

Challenges of Tail-based Trace Sampling

- Large feature dimension
 - Curse of dimensionality
 - Compromised performance and scalability
- Frequent pattern change
 - Microservices undergo continuous updates
 - New trace patterns could emerge

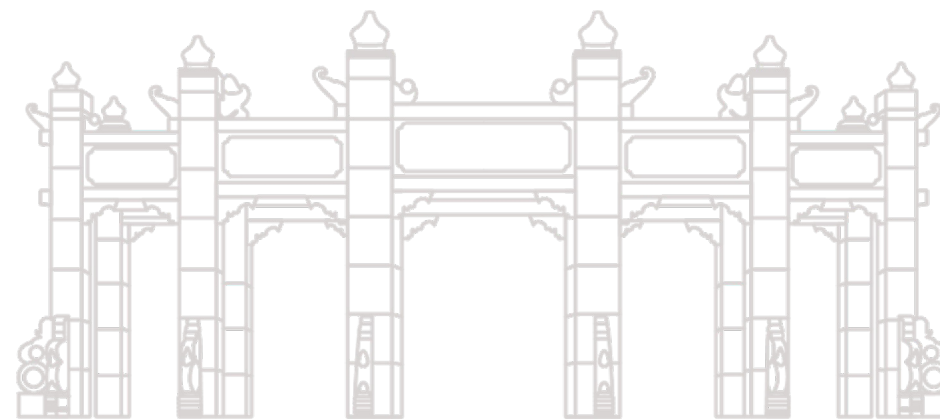


Microservices call graph collected at Uber

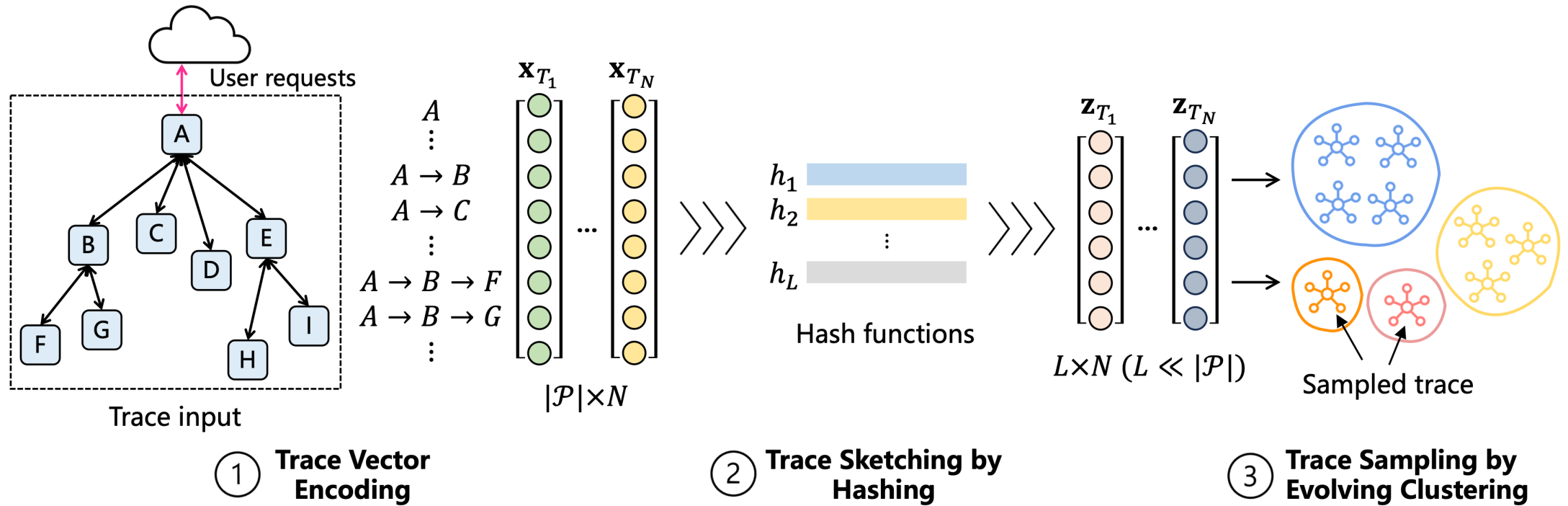
Not well addressed in existing approaches!



TraceMesh Overview & Design



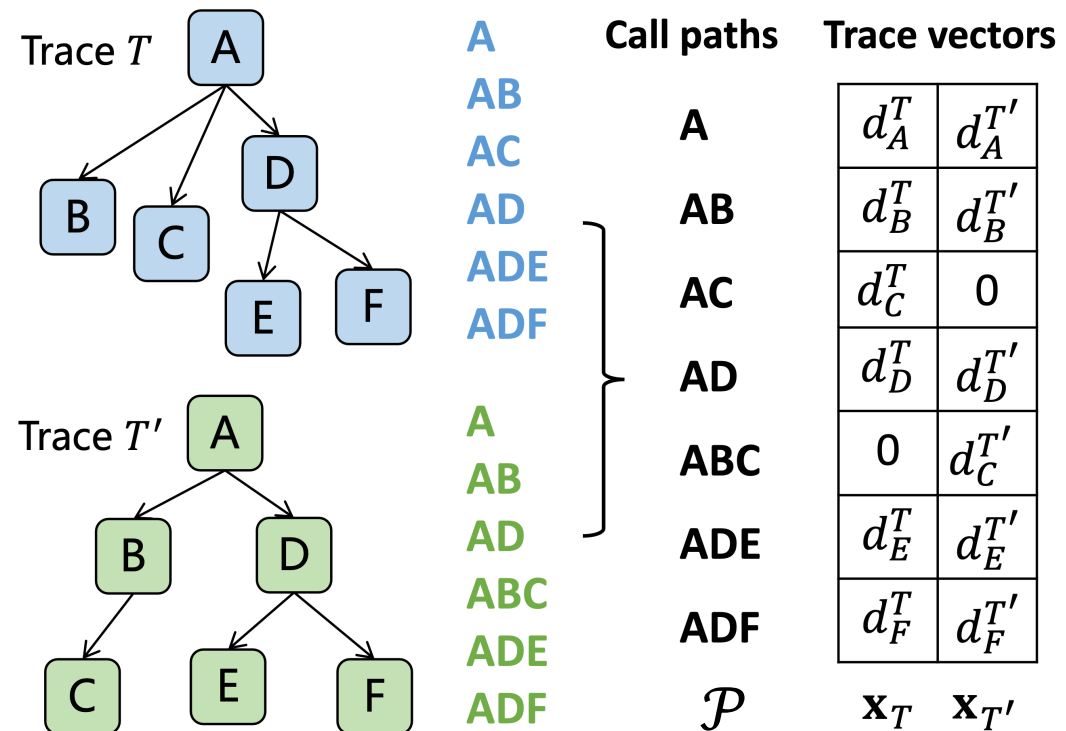
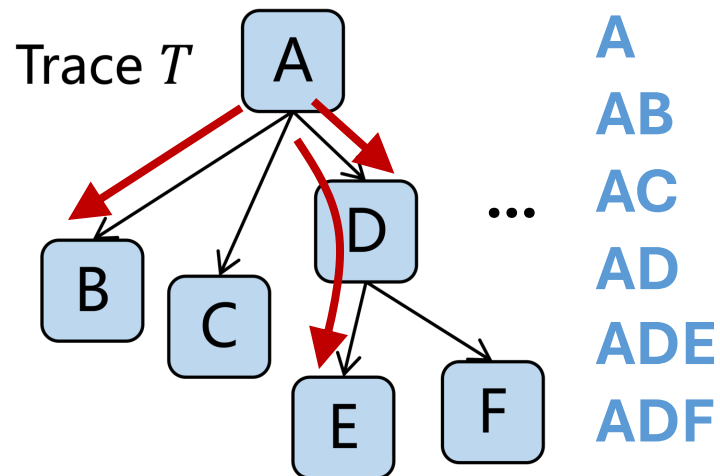
TraceMesh Overview



1. Transform graph traces into numerical feature vectors
2. Perform streaming trace vector encoding
3. Adaptively sample target traces by clustering

Step 1: Trace Vector Encoding

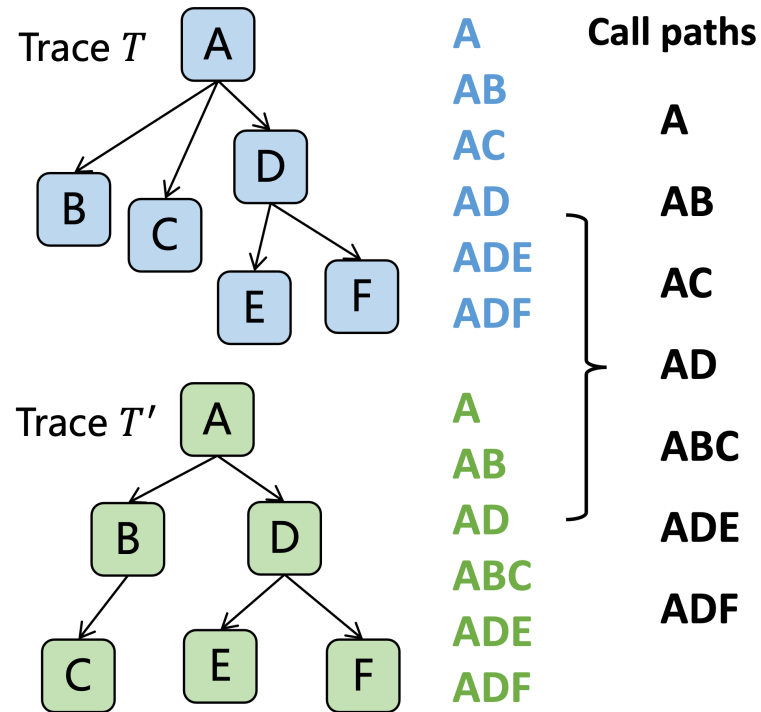
- Two types of trace features are extracted:
 - Structural feature
 - Temporal feature



The complete set of call paths can be large and dynamic!

Step 2: Trace Sketching by Hashing

- The first challenge:



Locality-Sensitive Hashing (LSH)

Project high-dimensional vectors into a low-dimensional space while preserving their similarity

$\mathcal{P}$ can be very large

Step 2: Trace Sketching by Hashing

- The projection process for a trace vector $\mathbf{x}$ in $\mathbb{R}^{|\mathcal{P}|}$:
 1. Instantiate L projection vectors $\mathbf{r}_1, \mathbf{r}_2, \dots, \mathbf{r}_L \in \{+1, -1\}^{|\mathcal{P}|}$
 2. Compute trace sketch vector as $\mathbf{z} = [h_{\mathbf{r}_1}(\mathbf{x}), \dots, h_{\mathbf{r}_L}(\mathbf{x})] \in \{+1, -1\}^L$:

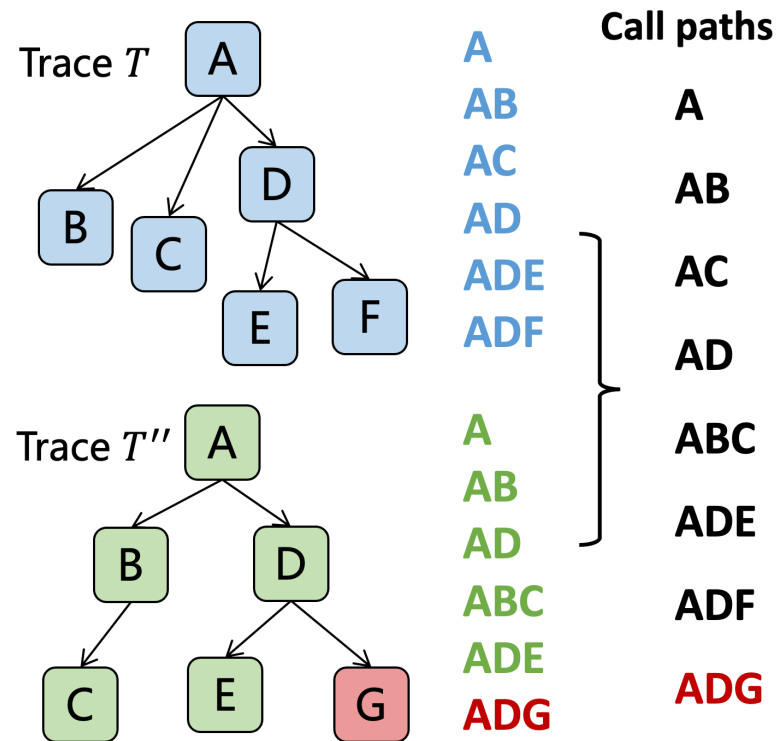
$$h_{\mathbf{r}_l}(x) = \begin{cases} +1 & \text{if } \mathbf{x} \cdot \mathbf{r}_l \geq 0 \\ -1 & \text{if } \mathbf{x} \cdot \mathbf{r}_l < 0 \end{cases} \quad \text{i.e., } h_{\mathbf{r}_l}(x) = \text{sign}(\mathbf{x} \cdot \mathbf{r}_l)$$

- Calculate the similarity for two traces by quantifying the agreement level between the sketch vectors:

$$\text{sim}(T, T') \propto \frac{|\{l: \mathbf{z}_T(l) = \mathbf{z}_{T'}(l)\}|}{L}$$

Step 2: Trace Sketching by Hashing

- The second challenge:



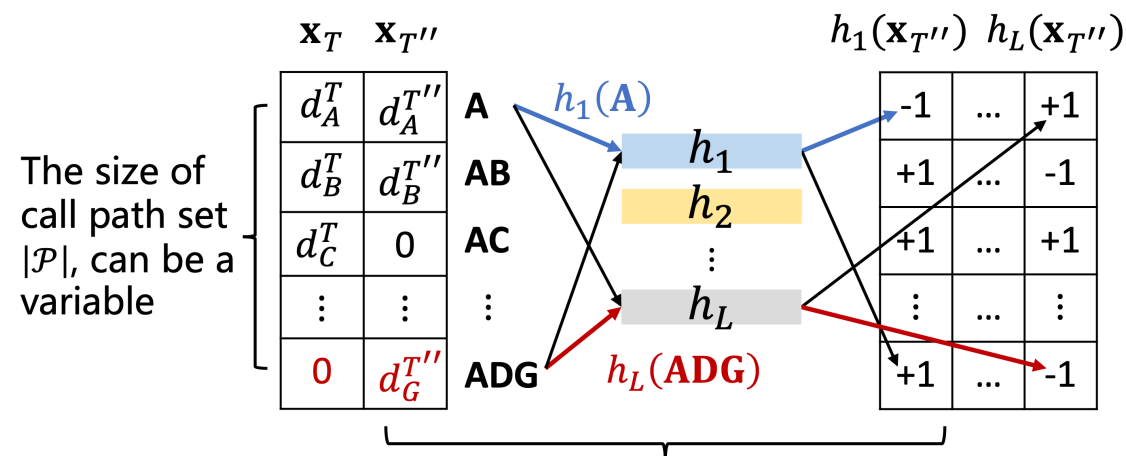
StreamHash

Replace project vectors with hash functions to deterministically map any call paths to $\{+1, -1\}$

$\mathcal{P}$ can be dynamic

Streaming Trace Vector Encoding

- L hash functions $h_1, h_2, \dots, h_L$ uniformly picked from a family $\mathcal{H}$ of hash functions



The trace T'' with a new path **ADG** can be seamlessly encoded as $\mathbf{z}_{T''}$

$$\mathbf{z}(l) = \text{sign}(\mathbf{y}(l)) = \text{sign}(\mathbf{x} \cdot h_l), l = 1, \dots, L$$

Fixed trace vector size L

$\mathbf{z}_{T''}(1)$	$\text{sign}(-1 * d_A^{T''} + \dots + (+1) * d_G^{T''})$
$\vdots$	$\vdots$
$\mathbf{z}_{T''}(L)$	$\text{sign}(+1 * d_A^{T''} + \dots + (-1) * d_G^{T''})$

Step 3: Trace Sampling by Evolving Clustering

- A **density-based clustering** approach for stream data
- Each **micro-cluster (MC)** has three attributes:
 - A weight w , determined by the traces inside it and time
 - A center c , the weighted center of the traces inside it
 - A radius r , the weighted avg distance from points to c
- Two types of MCs with different weight constraints:
 - **Potential MC (PMC)**: contains frequent and usual traces
 - **Outlier MC (OMC)**: contains potentially outlier traces

Step 3: Trace Sampling by Evolving Clustering

- When a new trace T arrives, TraceMesh
 - Merges it into the nearest PMC, and samples it with probability:

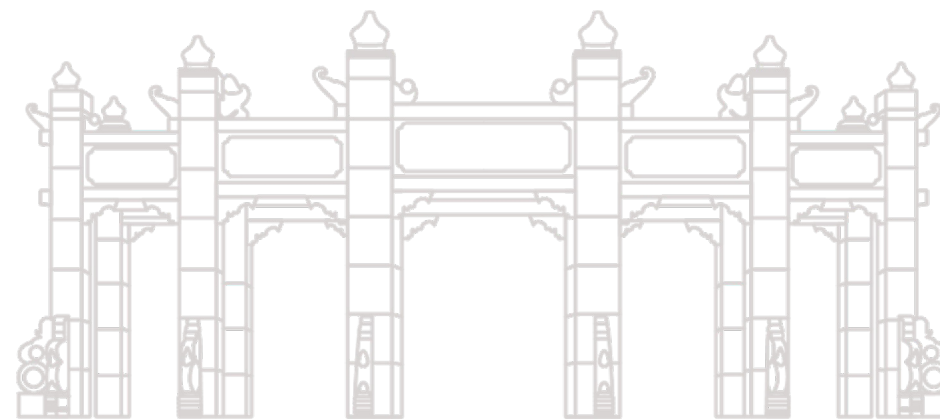
$$p_T = \mathcal{B} \times \left(1 - \frac{w_p^*}{\sum_{i=0}^{N_{pmc}} w^{(i)}}\right) \quad \mathcal{B}: \text{sampling budget}$$

- If failed, merges it into the nearest OMC, and samples it
- If failed, makes T a new OMC, and samples it

In this process, the **weight and role** of each MC will change.
TraceMesh **dynamically adjusts the sampling strategy** to identify the target traces without over-sampling recurring traces.



Evaluation



Datasets

- Traces from open-source microservices
 - Train Ticket
 - Online Boutique
- Traces from production systems

DATASET STATISTICS

Dataset	#Span	#Train	#Test	#Label
Train Ticket	201	653	31,814	624
Online Boutique	43	1,024	78,931	225
Industry	1,308	11,847	497,439	2,453

Metrics & Baselines

- Metrics

$$\text{Coverage} = \frac{\# \text{ target traces sampled}}{\# \text{ target traces}}$$

$$\text{Sampling rate} = \frac{\# \text{ traces sampled}}{\# \text{ traces}}$$

- Baselines

- Uniform sampling
- Perch [SoCC 2018]
- Sifter [SoCC 2019]
- Sieve [ICWS 2021]
- SampleHST [NOMS 2023]

Experimental Results

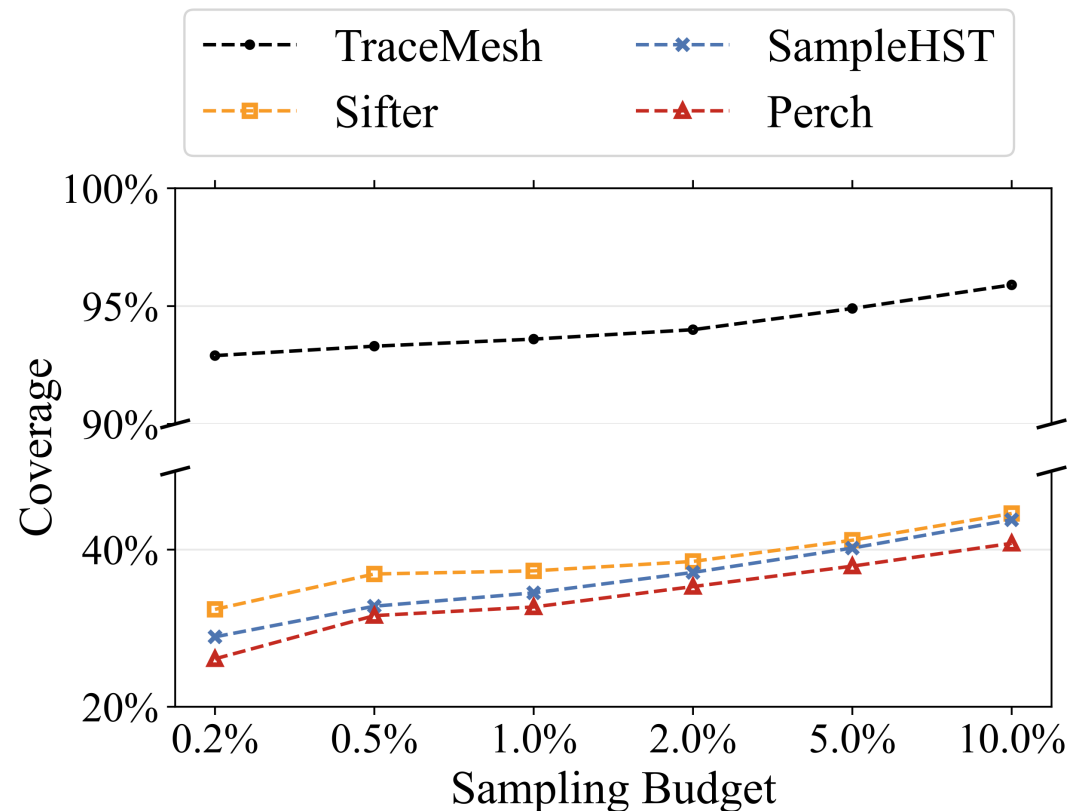
- TraceMesh achieves the best coverage
- TraceMesh has a relatively good sampling rate

EXPERIMENTAL RESULTS OF TRACE SAMPLING

Method	Train Ticket		Online Boutique		Industry	
	Coverage	Sampling Rate	Coverage	Sampling Rate	Coverage	Sampling Rate
Uniform	1.1%	1.00%	1.1%	1.00%	1.0%	1.00%
Sieve	71.8%	6.24%	82.2%	3.91%	61.1%	3.29%
Sifter	50.6%	1.28%	42.2%	1.13%	37.3%	1.16%
SampleHST	46.2%	1.19%	38.2%	1.07%	34.5%	1.10%
Perch	42.9%	1.00%	35.1%	0.98%	32.7%	0.95%
TRACEMESH	98.7%	2.34%	100%	1.35%	93.6%	0.83%

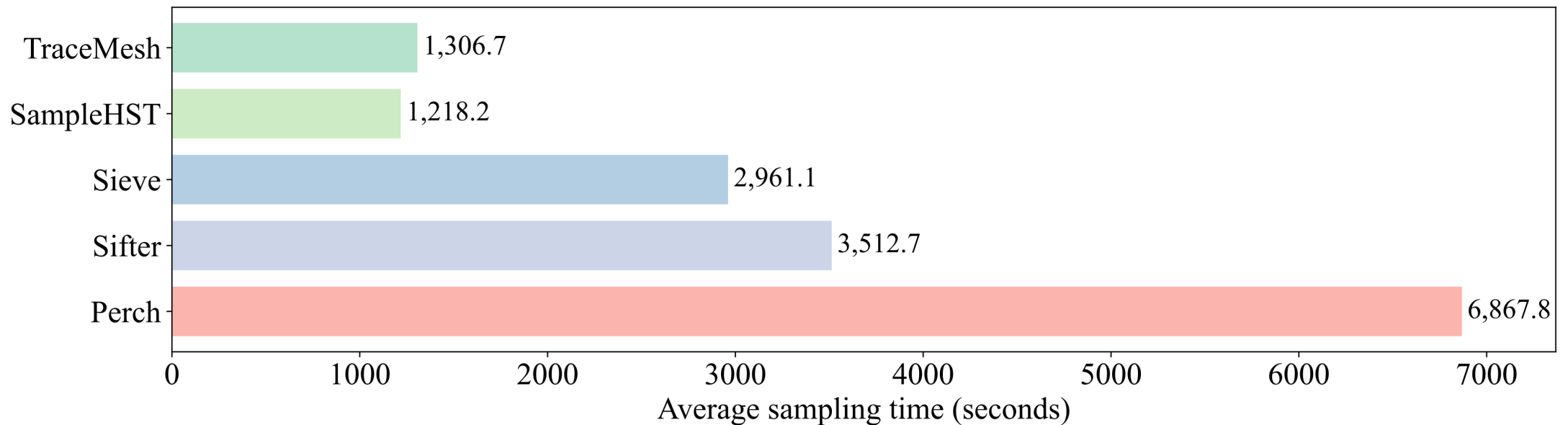
Coverage vs. Sampling Budget

- Coverage increases with larger sampling budget
- TraceMesh performs trace sampling more accurately



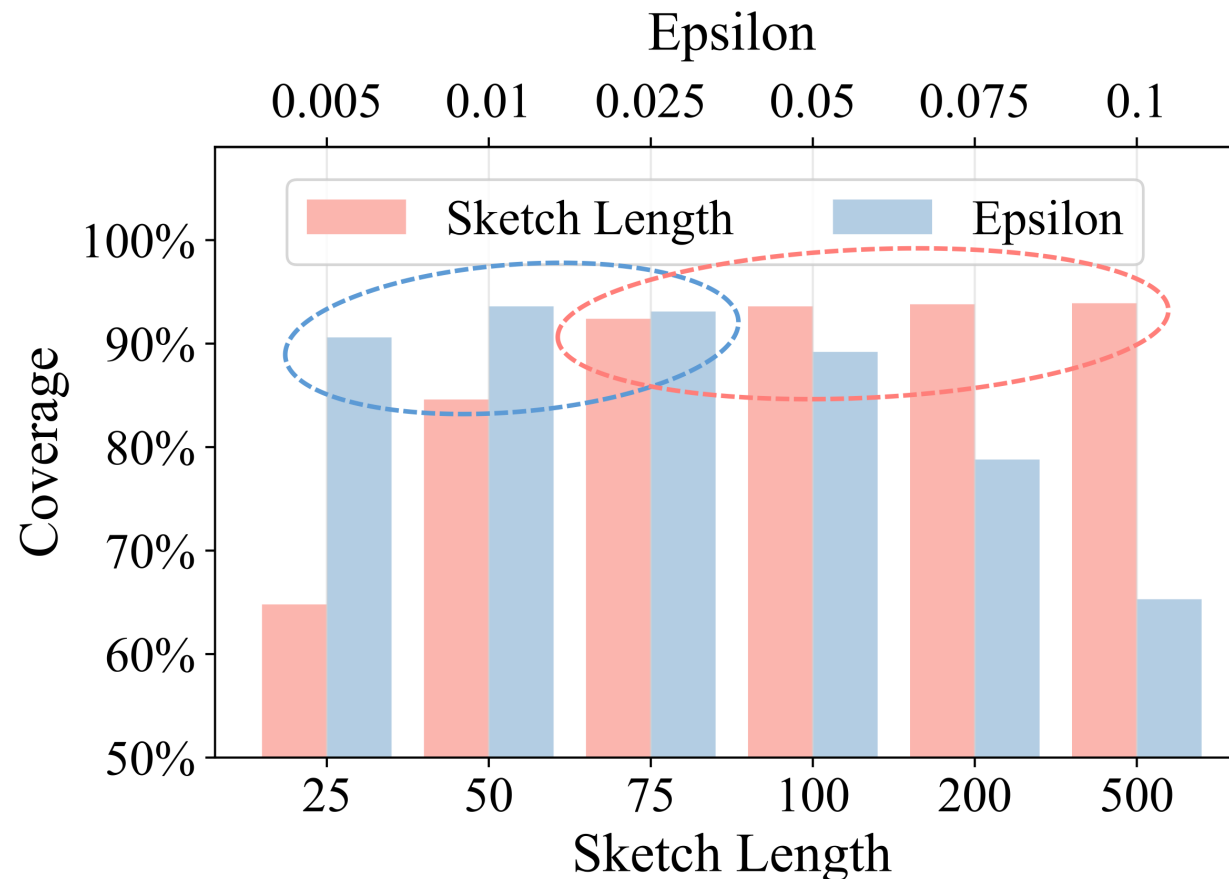
Efficiency

- TraceMesh performs trace sampling efficiently



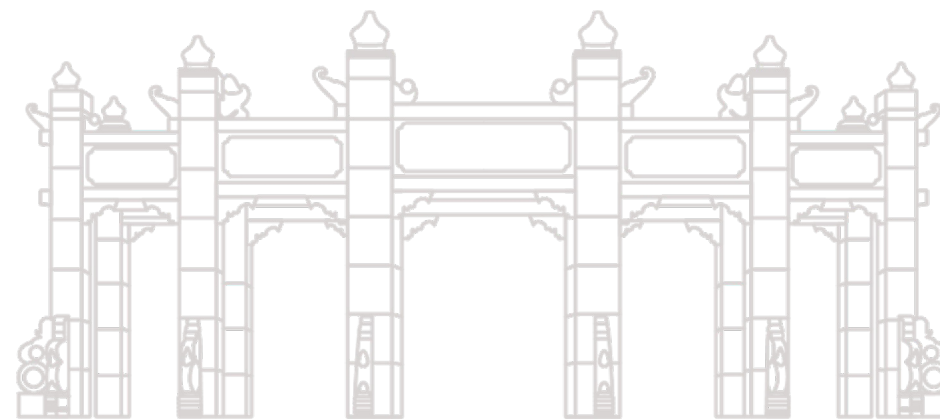
Parameter Sensitivity

- TraceMesh can achieve stable performance





Conclusion



Conclusion

- TraceMesh, a **scalable** and **streaming** trace sampler
 - Large feature dimension: LSH for dimension reduction
 - Frequent pattern changes: StreamHash and evolving clustering
- TraceMesh outperforms baselines in both sampling accuracy and efficiency



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Thanks!
Q & A

